IN THE CLAIMS

Please amend the claims as follows:

1. (Currently Amended) A program storage device, readablye by a machine, tangibly embodying programming instructions to perform method steps for constructing a call graph as a representation of a program, the programming instructions comprising:

selecting a program P for constructing a call graph representation thereof; wherein the program P contains zero or more fields F_E and one or more methods M_{M_1} .

- wherein each method M₁ in M_M has a single body B;
- wherein for each method M_2 in M_M , the call graph representation includes a corresponding node;
- wherein the call graph representation includes zero or more edges corresponding to connections between two or more of nodes;

determining identifying for each method M in M_M , a set of zero or more types S_M of objects that maywhich occur therein method M;

determining identifying for each field \underline{F} in \underline{F}_F ; a set of \underline{zero} or \underline{more} types S_F of objects that may be stored therein-field F;

determining identifying one or more the allocation sites inside the body <u>B</u> of <u>each</u> of method M;

determining the \underline{a} set of directly called methods M' inside the body \underline{B} of \underline{each} method M; and

determining thea set of virtually called methods M" inside the body B of each method M.

2. (Currently Amended) The program storage device according to claim 1, further comprising the programming instructions of:

adding T to the types S_M for each allocation of type T that occurs in the method M.

3. (Currently Amended) The program storage device according to claim 2, further comprising the programming instructions of:

for each direct call to <u>the methods</u> M' in a<u>the body B</u> of <u>the method M performing</u> the steps of:

adding any type that occurs in <u>the types</u> S_M and that is a subtype of <u>athe</u> type of a parameter of <u>the methods</u> M' to <u>types</u> $S_{M'}$; and

adding any type that occurs in <u>the types</u> $S_{M'}$ and that is a subtype of <u>the a</u> return type of <u>the methods</u> M' to <u>the types</u> S_{M} .

4. (Currently Amended) The program storage device according to claim 3, further comprising the programming instructions of:

for each virtual call to the methods M' in the body B of the method M:

using the types S_M , determine each of the methods M" that may be reached by thea dynamic dispatch:

adding any type that occurs in <u>the types</u> S_M and that is a subtype of the type of a parameter of <u>the methods</u> M'' to <u>a set</u> $S_{M''}$;

adding any type that occurs in the set $S_{M''}$ and that is a subtype of the return type of the methods M'' to the types S_{M} .

5. (Currently Amended) The program storage device according to claim 4, further the programming instructions of:

for each field F read by <u>the method M</u>, add any type that occurs in <u>the types S_F to the types S_M ; and</u>

for each field F with <u>the type T</u> written by <u>the method M</u>, add any type that occurs in the types S_M and that is a subtype of <u>the type T</u> to <u>the types S_F</u>.

6. (Original) The program storage device according to claim 1, further comprising the programming instructions of:

using the call graph computed above in a compiler as a basis for performing

optimizations such as inlining.

7. (Original) The program storage device according to claim 1, further comprising the programming instructions of:

using the call graph computed above in a reporting tool to report call graph information to a user.

8. (Currently Amended) A program storage device, readablye by a machine, tangibly embodying instructions to perform method steps for constructing a call graph as a representation of a program, the method comprising:

selecting a program P for constructing a call graph representation thereof; wherein the program P contains zero or more fields F_F and one or more methods M_{M_*}

- wherein each method M₁ in M_M has a single body B;
- wherein for each method M_2 in M_M , the call graph representation includes a corresponding node;
- wherein the call graph representation includes zero or more edges corresponding to connections between two or more of nodes;

determining identifying for each method M in M_M , a only one set of zero or more types S_M of objects that may which occur there in method M;

determining identifying for each field \underline{F} in \underline{F}_E , \underline{a} only one set of \underline{zero} or \underline{more} types S_E of objects that may be stored therein field \underline{F} ;

determining identifying one or more the allocation sites inside the body <u>B</u> of <u>each</u> of method M;

<u>determining a set of directly called methods M' inside the body B of each method</u>

<u>M; and</u>

determining a set of virtually called methods M" inside the body B of each method M.

9. (Currently Amended) The program storage device according to claim 8, further comprising the steps of:

determining the set of directly called methods M' inside the body <u>B</u> of <u>the method</u> M; and

determining the set of virtually called methods M" inside the body of method M. adding T to the types S_M for each allocation of type T that occurs in the method M.

10. (Currently Amended) A method for constructing a call graph <u>as a representation of a program</u>, the method comprising:

selecting a program P for constructing a call graph representation thereof;
wherein the program P contains zero or more fields F_F and one or more methods

M_M,

- wherein each method M₁ in M_M has a single body B;
- wherein for each method M_2 in M_M , the call graph representation includes a corresponding node;
- wherein the call graph representation includes zero or more edges corresponding to connections between two or more of nodes;

determining identifying for each method M in M_M , a set of zero or more types S_M of objects that may occur therein method M;

determining identifying for each field \underline{F} in \underline{F}_E , a set of \underline{zero} or \underline{more} types S_F of objects that may be stored therein field F;

determining identifying one or more the allocation sites inside the body <u>B</u> of <u>each</u> of method M;

determining thea set of directly called methods M' inside the body B of each method M; and

determining thea set of virtually called methods M" inside the body B of each method M.

- 11. (Currently Amended) The method according to claim 10, further comprising: adding T to the types S_M for each allocation of type T that occurs in the method M.
- 12. (Currently Amended) The method according to claim 11, further comprising:

 for each direct call to the methods M' in athe body B of the method M performing the steps of:

adding any type that occurs in <u>the types</u> S_M and that is a subtype of <u>thea</u> type of a parameter of <u>the methods</u> M' to <u>types</u> $S_{M'}$; and

adding any type that occurs in <u>the types</u> $S_{M'}$ and that is a subtype of <u>the a</u> return type of <u>the methods</u> M' to <u>the types</u> $S_{M'}$.

13. (Currently Amended) The method according to claim 12, further comprising: for each virtual call to the_methods M' in the body <a href="mailto:Body B] of the_method M:

using the types S_M , determine each of the methods M'' that may be reached by thea dynamic dispatch:

adding any type that occurs in <u>the types</u> S_M and that is a subtype of <u>thea</u> type of a parameter of <u>the methods</u> M'' to <u>a set</u> $S_{M''}$;

adding any type that occurs in <u>the set</u> $S_{M''}$ and that is a subtype of the return type of <u>the methods</u> M'' to <u>the types</u> S_M .

14. (Currently Amended) The method according to claim 13, further comprising:

for each field F read by \underline{the} method M, add any type that occurs in \underline{the} types S_{M} ; and

for each field F with <u>the</u> type T written by <u>the</u> method M, add any type that occurs in <u>the types</u> S_M and that is a subtype of <u>the type</u> T to <u>the types</u> S_F .

- 15. (Original) The method according to claim 10, further comprising the step of:
 using the call graph computed above in a compiler as a basis for performing optimizations such as inlining.
- 16. (Original) The method according to claim 10, further comprising the step of:
 using the call graph computed above in a reporting tool to report call graph
 information to a user.
- 17. (Cancelled) A method for constructing a scalable call graph, the method comprising: determining for each method M, only one set of types S_M of objects that may occur in method M; and determining for each field F, only one set of types S_F of objects that may be stored in field F; and determining the allocation sites inside the body of method M;
- 18. (Cancelled) The method to claim 17, further comprising the steps of: determining the set of directly called methods M' inside the body of method M; and determining the set of virtually called methods M" inside the body of method M.